

CONVERGENT EVOLUTION

Arriving at C++'s Concepts from a different path.

THERE'S NO GREAT INSIGHT HERE

- What I'm about to present is well-known information...
 - ...if you've read a lot of papers on the subject.
 - If you haven't read any papers, you might think that concepts are just a way to simplify template error messages.
- But the ideas here are *cool*!
 - I want to disseminate them to show you how cool they are.
- May help you reason about your programs in a new way.
- Also demonstrates that sometimes there really is only a single right answer when designing things.

PROGRAMS ARE PROOFS

- Called the “Curry-Howard correspondence”.
 - Types are propositions.
 - An object with that type proves the proposition that that type represents.
- Not 100% true in any real life programming system, *especially* not C++.
 - Evaluation scheme, side effects, cosmic rays, etc.
- But it’s fun to think that way sometimes.
 - Productive? I don’t know.
- **This idea ties all programming languages together.**

WHAT DOES IT MEAN?

- Functions are implications (IF):
 - Haskell: `A -> B`
 - `B (*)(A)`
 - `std::function<B(A)>`, I guess?
- Aggregates are conjunctions (AND):
 - `struct { A a; B b; }`
 - `std::tuple<A, B>`
 - *Many* ways in a lot of languages.
- Sum types are disjunctions (OR):
 - `union {A a; B b}`
 - Haskell: `data Foo = FooA A | FooB B`
- Parametric polymorphism is “for all” (\forall), also known as “universal types”:
 - Haskell: `id :: forall a. a -> a`
 - C++: `template<typename A> A id(A);`
- Existential types (\exists) are complicated.
 - TaPL (Pierce) says they’re pairs of values and types.
 - ML modules?
 - Encodable with universals anyway.

WHAT IS A CONSTRAINT?

- Say you want to sort a set. You have to prove that the elements of that set have a total order.
- `void qsort(void *base, size_t num, size_t size, int (*compar)(const void *, const void *))`
 - “compar” here is the *proof* that the order exists.
 - It is a *constraint* on the elements of the set.
 - Yes, this is a *very* rough approximation.
- **A constraint is an *extra parameter*.**
- They’re just lemmas in your proof.

SO... WHY CONCEPTS/TYPE CLASSES?

- Convenience!
- If it's an *established fact* that objects of type T have total order, then the sorting function shouldn't require me to explicitly supply that proof every time I want to sort a set of T objects.
- You're telling a function to JFGI.
 - Just f***ing get it from the environment!
 - Typeclasses in Haskell and Scala are actually *implicit parameters* indexed by *types* (yes, typeclasses are types too).
- Concepts and typeclasses allow you to state facts about types in a single place and let those facts be known everywhere, just by mentioning the constraint.

DRINK THE KOOL-AID

Haskell

```
sort :: Ord a => [a] -> [a]
```

```
class Eq a => Ord a where  
    (<=) :: a -> a -> Bool
```

```
class Eq a where  
    (==) :: a -> a -> Bool
```

C++

```
template<typename T> requires Ord<T>  
list<T> sort(list<T>);
```

```
template<typename T>  
concept bool Ord = requires (T a, T b)  
    { {a <= b} -> bool; ... }
```

```
template<typename T>  
concept bool Eq = requires (T a, T b)  
    { {a == b} -> bool; ... }
```

FURTHER READING

- Oliveira, Moors, Odersky: Type Classes as Objects and Implicits. 2010.
- Bernardy, Jansson, Zalewsky, Schupp: Generic Programming with C++ Concepts and Haskell Type Classes – A Comparison. 2010.